

Programming languages and particle physics

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Workshop on Analysis Description Languages for the LHC 6-8 May 2019, Fermilab LPC



https://indico.cern.ch/event/769263/

An analysis description language (ADL) is a human readable declarative language that unambiguously describes the contents of an analysis in a standard way, independent of any computing framework.

Adopting ADLs would bring numerous benefits for the LHC experimental and phenomenological communities, ranging from analysis preservation beyond the lifetimes of experiments or analysis software to facilitating the abstraction, design, visualization, validation, combination, reproduction, interpretation and overall communication of the



But that just ended a few minutes ago.

(This talk is not a summary of the workshop; come to tomorrow's LPC Physics Forum at 1:30pm.)



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(This talk is not a summary of the workshop; come to tomorrow's LPC Physics Forum at 1:30pm.)

Instead, let's take a step back...

You cannot step into the same river twice.

Heraclitus



Because, you know, it's different water.



So why do we say it's the same river?



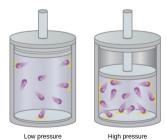
The river is an abstraction.

We associate an enormous number of microscopic states ("molecules here, molecules there") with a single macroscopic state ("the river").



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It's an abstraction like thermodynamics; it can be exact with the right definitions.



```
double bessel i0(double x) {
    double out;
    if (fabs(x) < 8.0) {
        double y = x \star x;
        double ans1 = 57568490574.0 + y*(-13362590354.0 + y*(651619640.7
                      + y*(-11214424.18 + y*(77392.33017 + y*(-184.9052456)))));
        double ans2 = 57568490411.0 + y*(1029532985.0 + y*(9494680.718
                      + v*(59272.64853 + v*(267.8532712 + v*1.0)));
        out = ans1 / ans2;
    else ·
        double z = 8.0 / fabs(x):
        double v = z * z;
        double xx = fabs(x) - 0.785398164;
        double ans1 = 1.0 + y* (-0.1098628627e-2 + y* (0.2734510407e-4
                      + v*(-0.2073370639e-5 + v*0.2093887211e-6)));
        double ans2 = -0.1562499995e-1 + v*(0.1430488765e-3
                      + v*(-0.6911147651e-5 + v*(0.7621095161e-6
                      - y*0.934935152e-7)));
        out = sqrt(0.636619772/fabs(x))*(cos(xx)*ans1 - z*sin(xx)*ans2);
    return out;
```

Most of computer science is about abstracting details, too.



```
\leftarrow one value goes in
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        out = sqrt(0.636619772/fabs(x))*(cos(xx)*ans1 - z*sin(xx)*ans2);
                                    \leftarrow one value comes out
    return out;
```





The abstraction is cumulative:

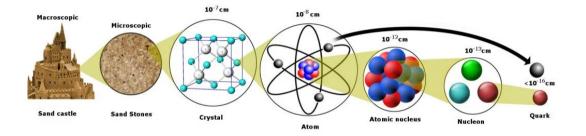
Every function/class/module has an interior and an interface—minimizing

 $\frac{\#\text{external parameters}}{\#\text{internal parameters}}$

reduces the mental burden on programmers and users.

Science has layers of abstraction

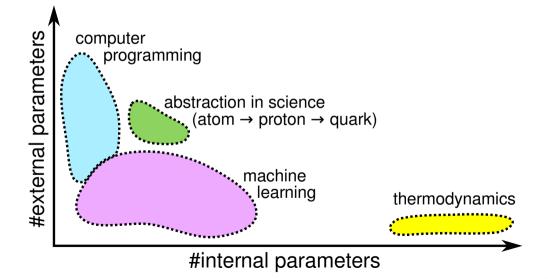




These are approximate, taking advantage of a separation of scales.

(cartoon diagram, not to scale)

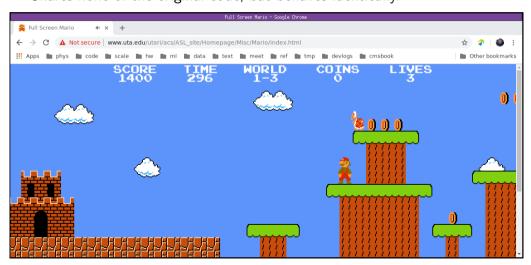




Software interfaces can be exact, despite radical internal differences.

R

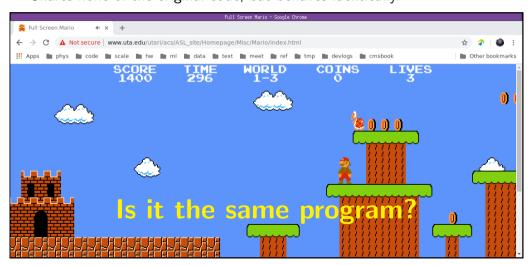
Super Mario Bros. entirely rewritten in Javascript by Josh Goldberg.
 Shares none of the original code, but behaves identically.



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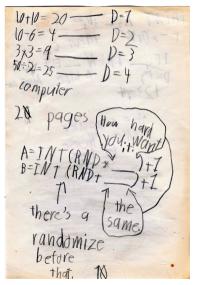
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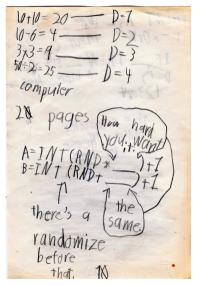
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Which meant Pascal. Pascal was "real," and BASIC was not.

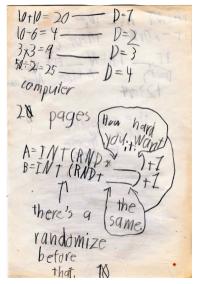




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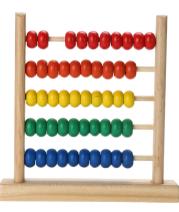
Which meant Pascal. Pascal was "real," and BASIC was not.

But ultimately, not even assembly code is real in the sense that I'm meaning here.



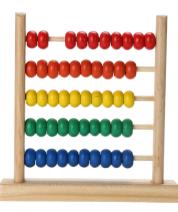


The objectively real part of a computer is a set of physical states.



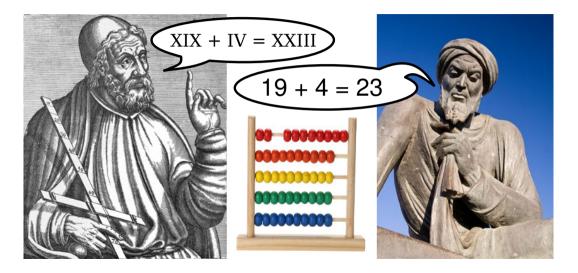


The objectively real part of a computer is a set of physical states *that we interpret* as computations.



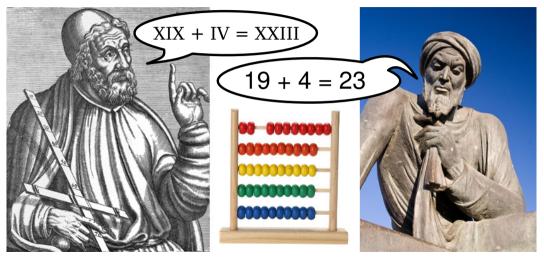
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(And some languages are better at it than others.)



Programming languages differ in their degree of abstraction, but *all* programming languages are for humans, not computers.



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Each one re-expresses the programmer's intent in terms of another:

CMSSW configuration Python runtime C source code machine instructions logic gates

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Only the last level actually pushes the abacus beads.



Ada of Lovelace's algorithm for computing Bernoulli numbers was written for a computer that never ended up being invented, but it was a program.

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John McCarthy, creator of Lisp: "This EVAL was written and published in the paper and Steve Russel said, 'Look, why don't I program this EVAL?' and I said to him, 'Ho, ho, you're confusing theory with practice—this EVAL is intended for reading, not for computing!' But he went ahead and did it."



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APL (ancestor of MATLAB, R, and Numpy) was also a notation for describing programs years before it was executable. The book was named *A Programming Language*.



Programmers had to manually translate these notations into instruction codes.

That's why it was called "coding."



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That's why it was called "coding."

Von Neumann called assembly language "a waste of a valuable scientific computing instrument—using it for clerical work!"



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Now that our programming languages *do* push abacus beads, software engineering has become an odd discipline: saying something is the same as making it.

And yet, we still get it wrong.





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But what about speed? Don't we choose languages for speed?



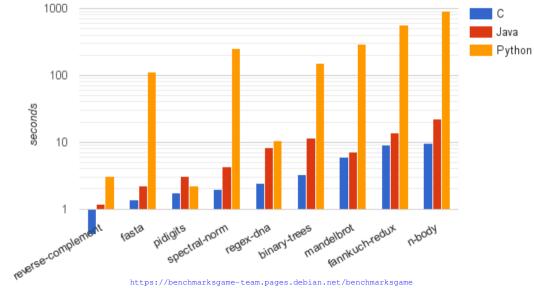
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But what about speed? Don't we choose languages for speed?

"There's no such thing as a 'fast' or 'slow' language." — so sayeth the StackOverflow

Except Python. Python is slow, right?





https://benchmarksgame-team.pages.debian.net/benchmarksgame

But it really isn't the language; it's the implementation.



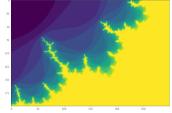


```
def run(height, width, maxiterations=20):
    y, x = numpy.ogrid[-1:0:height*1j, -1.5:0:width*1j]
    c = x + y + 1 i
    fractal = numpy.full(c.shape, maxiterations,
                                   dtype=numpy.int32)
    for h in range(height):
        for w in range(width):
                                                  # for each pixel (h, w)...
            z = c[h, w]
            for i in range(maxiterations):
                                                  # iterate at most 20 times
                z = z * * 2 + c[h, w]
                                                  # applying z \rightarrow z^2 + c
                if abs(z) > 2:
                                                  # if it diverges (|z| > 2)
                    fractal[h, w] = i
                                                  # color the plane with the iteration number
                    break
                                                  # we're done, no need to keep iterating
```

return fractal

But it really isn't the language: it's the implementation.





```
import numpy, numba
@numba.jit
def run(height, width, maxiterations=20):
    y, x = numpy.ogrid[-1:0:height*1j, -1.5:0:width*1j]
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```

return fractal

Now 50 \times faster, about equal to C code (-03).



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Same language (subset), completely different implementation.



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Same language (subset), completely different implementation.



Pure Python is slower than Numba or C because it has more hurdles in the way: dynamic typing, pointer-chasing, garbage collection, hashtables, string equality...

Greg Owen's talk on Spark 2.0



Performance Improvements in Spark 2.0

Greg Owen 2016-05-25





Volcano Iterator Model

```
Standard for 30 years: almost all databases do it
```

```
Each operator is an "iterator"
that consumes records from
its input operator
```

```
class Filter {
  def next(): Boolean = {
    var found = false
    while (!found && child.next()) {
      found = predicate(child.fetch())
    }
    return found
}
def fetch(): InternalRow = {
```

child.fetch()

23

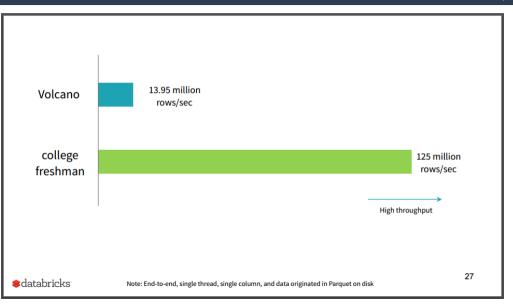


What if we hire a college freshman to implement this query in Java in 10 mins?

```
select count(*) from store_sales
where ss item sk = 1000
```

```
var count = 0
for (ss_item_sk in store_sales)
{
    if (ss_item_sk == 1000) {
        count += 1
     }
}
tdatabricks
```

25





How does a student beat 30 years of research?

Volcano

- 1. Many virtual function calls
- 2. Data in memory (or cache)
- 3. No loop unrolling, SIMD, pipelining

Hand-written code

- 1. No virtual function calls
- 2. Data in CPU registers
- 3. Compiler loop unrolling, SIMD, pipelining

Take advantage of all the information that is known after query compilation

databricks



So although it's the implementation, not the language, that's slow,

that implementation can be hampered by the flexibility that the language promises.

lukeplant.me.uk

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We need less powerful languages

by Luke Plant

Posted in: Python, Haskell, Django – November 14, 2015 at 11:46

Translations of this post (I can't vouch for their accuracy):

Japanese

Many systems boast of being 'powerful', and it sounds difficult to argue that this is a bad thing. Almost everyone who uses the word assumes that it is always a good thing.

The thesis of this post is that in many cases we need less powerful languages and systems.

Before I get going, there is very little original insight in this post. The train of thought behind it was set off by reading Hofstadter's book **Gödel**, **Escher**, **Bach** — **an Eternal Golden Braid** which helped me pull together various things in my own thinking where I've seen the principle in action. Philip Wadler's post on **the rule of least power** was also formative, and most of all I've also taken a lot from the content of **this video from a Scala conference about everything that is wrong with Scala**, which makes the following fairly central point:

Every increase in expressiveness brings an increased burden on all who care to understand the message.



Domain-specific languages:

specialized languages for narrowly defined problems.

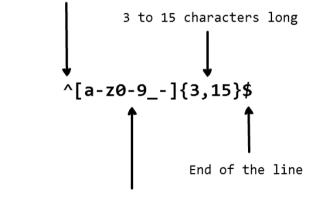
- Main purpose: reduces complexity, the mental clutter that obscures general-purpose languages.
- Secondary purpose: limited flexibility allows for streamlined implementations.



Any guesses?

Regular expressions

Start of the line



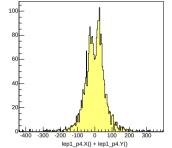
letters, numbers, underscores, hyphens

Domain-specific languages that you're probably already using





ttree->Draw("lep1_p4.X() + lep1_p4.Y()");



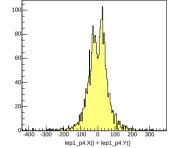
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ttree->Draw("lep1_p4.X() + lep1_p4.Y()");

Looping and reducing constructs:



for (int i0; i0 < 3; i0++) {
 for (int j2; j2 < 5; j2++) {
 for (int j3; j3 < 2; j3++) {
 for (int j3; j3 < 2; j3++) {
 int i1 = fResults[j2][j3];
 use the value of fMatrix[i0][i1]
 }
 }
 }
 }
}</pre>

 $\texttt{Length}(\cdot) \quad \texttt{Sum}(\cdot) \quad \texttt{Min}(\cdot) \quad \texttt{Max}(\cdot) \quad \texttt{MinIf}(\cdot, \cdot) \quad \texttt{MaxIf}(\cdot, \cdot) \quad \texttt{Alt}(\cdot, \cdot) \quad \texttt{MaxIf}(\cdot, \cdot) \quad \texttt{MaxI$



Makefiles

🔄 *makefile 🗙	
makefile:	<u>^</u>
all: hello	
clean:	
-rm main.o hello.exe hello	
hello: main.o	
g++ -g -o hello main.o	
main.o: main.cpp	
g++ -c -g main.cpp	
	-
	Þ



Format strings

printf/scanf: distinct syntax from C/C++, must be quoted
printf("Error 0x%04x: %s", id, errors[id]);

I/O streams: defined within $C/C{++}$ via operator overloading



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printf/scanf is "external" and I/O streams is "internal" (embedded)

External versus internal (embedded) domain-specific languages



External: SQL has a distinct syntax from Python; must be quoted in PySpark.

```
import pyspark
pyspark.sql("""
   SELECT CONCAT(first, " ", last) AS fullname, AVG(age)
        FROM my_table WHERE age BETWEEN 18 AND 24
        GROUP BY fullname
""")
```

Internal (embedded): SparkSQL is an equivalent language, defined within Python.

```
import pyspark.sql.functions as F
df = pyspark.read.load("my_table")
(df.withColumn("fullname",
            F.concat(F.col("first"), F.lit(" "), F.col("last")))
    .select("fullname", "age")
    .where(df.age.between(18, 24))
    .groupBy("fullname")
    .agg(F.mean("age")))
```



Objection: a collection of libraries and operator overloads isn't a language!



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My answer: programming languages are human modes of expression, implemented using other programming languages, all the way down.

What matters is whether it's a coherent set of concepts, not whether it was implemented by a parser.



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My answer: programming languages are human modes of expression, implemented using other programming languages, all the way down.

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(One might as well argue about the distinction between languages and dialects.)



Perhaps the most widespread domain-specific language in data analysis:

SQL



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SQL

But we rarely use it in particle physics. Why?



"Momentum of the track with $|\eta| < 2.4$ that has the most hits."

```
Track *best = NULL;
for (int i = 0; i < tracks.size(); i++) {</pre>
  if (fabs(tracks[i]->eta) < 2.4)</pre>
    if (best == NULL |
        tracks[i]->hits.size() > best->hits.size())
      best = tracks[i];
if (best != NULL)
  return best->pt;
else
  return 0.0;
```



"Momentum of the track with $|\eta| < 2.4$ that has the most hits."

```
WITH hit stats AS (
  SELECT hit.track id, COUNT(*) AS hit count FROM hit
    GROUP BY hit.track id),
 track sorted AS (
    SELECT track.*,
    ROW NUMBER() OVER (
     PARTITION BY track.event id
     ORDER BY hit stats.hit count DESC)
  track ordinal FROM track INNER JOIN hit stats
    ON hit stats.track id = track.id
    WHERE ABS(track.eta) < 2.4)
 SELECT * FROM event INNER JOIN track sorted
   ON track sorted.event id = event.id
WHERE
  track sorted.track ordinal = 1
```



The problem is that collisions produce a variable number of particles per event: the tables are "jagged."



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This can be described using SQL's relational concepts:

- separate tables for events and particles
- ▶ linked by a common "event number" index.

But each type of particle has to be a separate table and each operation has to be **INNER JOIN**ed to maintain events as objects.



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But each type of particle has to be a separate table and each operation has to be **INNER JOIN**ed to maintain events as objects.

SQL makes particle physics problems harder, not easier, which defeats the point.



Would a domain specific language for particle physics

- make analysis code easier to read?
- make mistakes more evident?
- make it easier to synchronize analyses from different groups/experiments?
- make it easier to preserve them in executable/recastable form?
- highlight physics concepts, like control regions, systematic variations, event weights, combinatorics with symmetries?
- hide irrelevant concepts like managing files, memory, load balancing, and other performance tweaks?



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That was the subject of the Analysis Description Language Workshop.

In fact, about that SQL...



and indexing (SZOMME) (rathienal) immutable? Data types Database = event Itable table tuble electrons jets tracky Select X, J, Z from electrons MENL

In fact, about that SQL...





Why hasn't this been done before?

(Why hasn't it succeeded before?)



I think the answer is cultural, so I'll take a historical perspective...



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Starting in 1880.

The U.S. Census's problem



The U.S. does a census every 10 years. The 1880 census took 8 years to process. \longrightarrow Big data problem!

The U.S. Census's problem



The U.S. does a census every 10 years. The 1880 census took 8 years to process. \longrightarrow Big data problem!

Held a competition for a new method; winner was 10 \times faster than the rest:

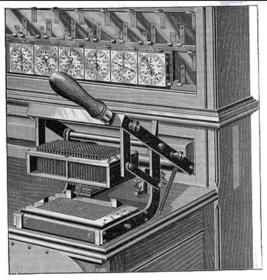


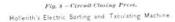


Census records on punch cards, which filtered electrical contacts









Wired to a machine that opens a door for each matching pattern



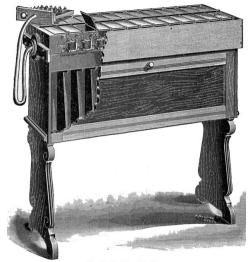
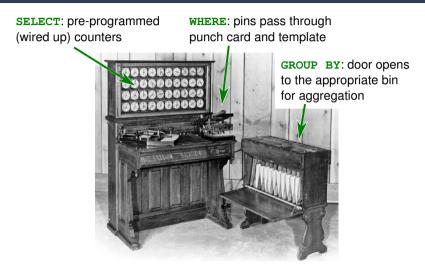


Fig. 3.-Sorting Machine. Hollerith's Electric Sorting and Tabulating Machine.

It was an SQL machine: 3 basic clauses of most SQL queries





SELECT name WHERE literate GROUP BY marital_status



Herman Hollerith (inventor) incorporated the Tabulating Machine Company, which after a series of mergers became International Business Machines (IBM) in 1924.



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Most recently as "map-reduce."





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2004: published as a paper by Jeffrey Dean and Sanjay Ghemawat. 2006: reimplemented as open-source software: Apache Hadoop.

Problems like "index all webpages" plug into this framework.

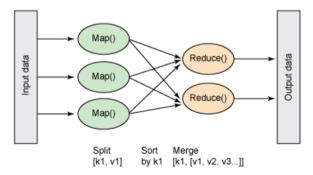


SELECT-WHERE: filter and transform each input to a $\langle key, value \rangle$ pair.

```
def map(webpage):
    for word in webpage.split():
        if not stopword(word):
            yield (word, webpage)
```

GROUP BY: collect and transform all values with a given key.

def reduce(word, webpages):
 index[word] = set()
 for webpage in webpages:
 index[word].add(webpage)





That's how statisticians encountered computing.

Physics encountered computing differently.









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Los Alamos group led by Nicholas Metropolis, developed Monte Carlo techniques for physics problems.



The actual programming was performed by these six women









Kathleen McNulty

Frances Bilas E

Betty Jean Jennings



Ruth

Lichterman



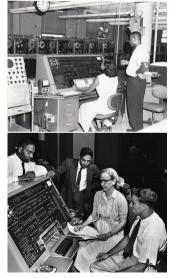
Elizabeth Snyder

Marlyn Wescoff

Eckert-Mauchly Computer Corporation \rightarrow Remington Rand



Mauchly and Eckert "went into industry" selling computers; the first one (UNIVAC) to the U.S. Census.

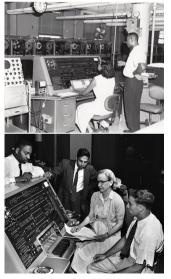




Mauchly and Eckert "went into industry" selling computers; the first one (UNIVAC) to the U.S. Census.

1950: Short Code, the first executable high-level language: a transliterated interpreter of mathematical formulas.

math: X3 = (X1 + Y1) / X1 * Y1
code: X3 03 09 X1 07 Y1 02 04 X1 Y1
50× slower than machine code because it was interpreted.



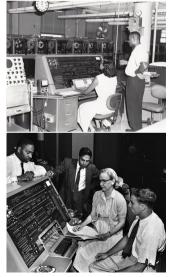
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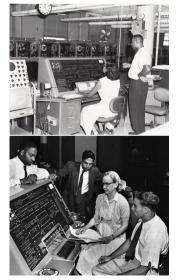
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Meanwhile, IBM developed FORTRAN: 1954-1957.







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"The whole concept of ZEBRA is a manifestation of one of FORTRAN 77's needs." — Bebo White in 1989







EUROPEAN ORGANIZATION FOR NUCLEAR RESEARCH

CERN-DD /89/ 18

May 16, 1989

The Comparison and Selection of Programming Languages for High Energy Physics Applications

Bebo White

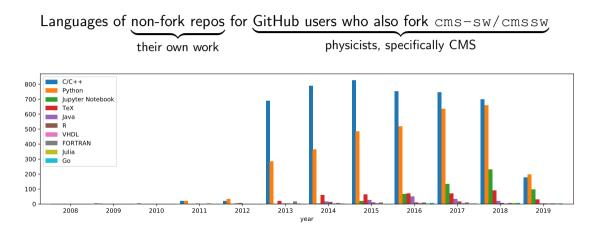
Data Handling Division, CERN and SLAC Computing Services



Zanella [32] has said " If HEP wishes to keep to its level of achievement, credibility and excellence, then it needs an injection of bright young computer-wise scientists and engineers." This means that HEP cannot become "an island." HEP applications must be able to utilize "state of the art" facilities in all areas of applicability including data processing. HEP must be able to take advantage of the technological advancements in other arenas of science and engineering. Many of these advancements are occurring in fields which are presently not software compatible with HEP. Much of the work being done in embedded systems with Ada or telecommunications with C could be of great interest and applicability in HEP computing environments. The unified physics computing environment anticipated for the 1990s should be able to take full advantage of these facilities and the physicists and engineers of the 1990s should be able to take full advantage of their unified physics computing environment.

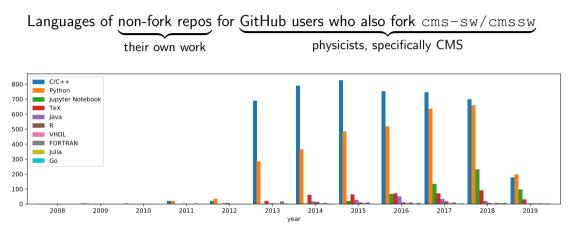
Are we halfway through the second major language shift?





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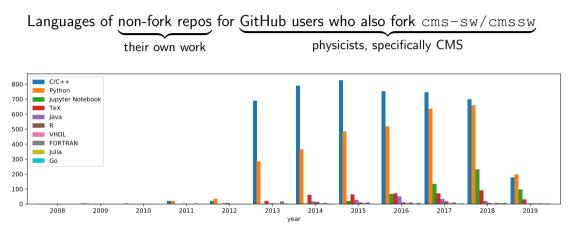




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Are we halfway through the second major language shift?





The shift from Fortran to C++ was a decision made by collaboration leaders.

What we see here are individuals choosing a language for their own work.

Workshop on Analysis Description Languages for the LHC 6-8 May 2019, Fermilab LPC



https://indico.cern.ch/event/769263/

An analysis description language (ADL) is a human readable declarative language that unambiguously describes the contents of an analysis in a standard way, independent of any computing framework.

Adopting ADLs would bring numerous benefits for the LHC experimental and phenomenological communities, ranging from analysis preservation beyond the lifetimes of experiments or analysis software to facilitating the abstraction, design, visualization, validation, combination, reproduction, interpretation and overall communication of the



Informal summary of the workshop at tomorrow's LPC Physics Forum at 1:30pm.